Logic in Computer Science

Lecture 07
Hoare Logic and Program Verification

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Verification methods may be classified according to the following main criteria:

- ***Proof-based vs. model-based*** - if a soundness and completeness theorem holds, than:
  - proof = valid formula = true in all models;
  - model-based = check satisfiability in one model

- ***Degree of automation*** - fully automated, partially automated, or manual

- ***Full- vs. property-verification*** - a single property vs. full behavior

- ***Domain of application*** - hardware or software; sequential or concurrent; reactive or terminating; etc.

- ***Pre- vs. post-development***

Used to verify sequential programs with infi nite state and complex data.

- Proof based

- Semi-automatic —some steps cannot be carried out algorithmically by a computer.

- Property-oriented

- Application domain:Sequential, transformational programs

- Pre/post development: the methods can be used during the development process to create small proofs that can be subsequently combined into proofs of larger program fragments.

Why should we specify and verify code

- A formal specification is less ambiguous.

- Experience has shown that verifying programs w.r.t. formal specifications can significantly cut down the duration of software development and maintainance by eliminating most errors in the planning phase.

- Makes debugging easier

- Software built from formal specifications is easier to reuse.

- Verification of safety-critical software ***guarantees*** safety; testing does not.

- Many examples of software-related catastrophies due to lack of verification.

  – Arianne rocket exploded immediately after launch

  – Lost control of Martian probe

  – Y2K problem

A Framework for Software Verification

As a software developer, you may get an order from a customer, which provides an informal description of your task.

- Convert the informal description D of an application domain into an "equivalent" formula Φ_D of some symbolic logic.

- Write a program P which is meant to realize Φ_D in the programming environment required by the customer.

- Prove that P satisfi es Φ_D.

We use a language with simple integer and boolean expressions, and simple commands: assignment, if, and while commands.

$$
\begin{aligned}
E \ &::= \ n \,|\, x \,|\, (-E) \,|\, (E+E) \,|\, (E-E) \,|\, (E*E) \\
B \ &::= \ \texttt{true} \,|\, \texttt{false} \,|\, (!B) \,|\, (B\&B) \,|\, (B||B) \,|\, E < E \\
C \ &::= \ x = E \,|\, C;C \,|\, \texttt{if } B \,\{C\} \texttt{ else } \{C\} \,|\, \texttt{while } B \,\{C\}
\end{aligned}
$$

Example:

```
y = 1;
z = 0;
while (z != x) {
    z = z + 1;
    y = y * z;
}
```

We need to be able to express the following statement: "If the execution of a program fragment P starts in a state satisfying Φ, then the execution of P ends in a state satisfying Ψ. We denote this by:

$$(\!|\Phi|\!) \; P \; (\!|\Psi|\!)$$

and we call this construct a ***Hoare triple***. Φ is called the ***precondition***, and Ψ is called the ***postcondition***.

Example: Assume that the specifi cation of a program P is ***"to calculate a number whose square is less that x."*** Then, the following assertion should hold:

$$(\!|x > 0|\!) \; P \; (\!|y \cdot y < x|\!)$$

It means: if we start execution in a state where $x > 0$, then the execution of P ends with a state where $y^2 < x$.

What happens if the execution starts with $x \leq 0$? We don't know!

Both these examples realize the specification $(\!|x > 0|\!)\ P\ (\!|y \cdot y < x|\!)$.

$$(\!|x > 0|\!)$$
$$\mathtt{y} = \mathtt{0}$$
$$(\!|y \cdot y < x|\!)$$

$$(\!|x > 0|\!)$$
```
y = 0
while (y * y < x) {
    y = y + 1
}
y = y - 1
```
$$(\!|y \cdot y < x|\!)$$

- ***Partial correctness:*** we ***do not require*** the program to terminate.

- ***Total correctness:*** we ***do require*** the program to terminate.

Definition (partial correctness): We say that the triple $(\!|\Phi|\!)\ P\ (\!|\Psi|\!)$ is satisfied under partial correctness if, for all states which satisfy Φ, the state resulting from P's execution satisfies the postcondition Ψ, provided that P actually terminates. In this case we write

$$\models_{par} (\!|\Phi|\!)\ P\ (\!|\Psi|\!)$$

Definition (total correctness): We say that the triple $(\!|\Phi|\!)\ P\ (\!|\Psi|\!)$ is satisfied total partial correctness if, for all states in which P is executed and which satisfy the precondition Φ, P is guaranteed to terminate, and the state resulting from P's execution satisfies the postcondition Ψ. In this case we write

$$\models_{tot} (\!|\Phi|\!)\ P\ (\!|\Psi|\!)$$

The following statement

$$\models_{par} (\!|\Phi|\!)\ \texttt{while true \{ x = 0; \}}\ (\!|\Psi|\!)$$

holds for all Φ and Ψ. The corresponding total correctness statement does not hold.

`Succ:`

```
a = x + 1;
if (a - 1 == 0 {
    y = 1;
} else {
    y = a;
}
```

We have:

$$\models_{par} (\!|\top|\!)\ \texttt{Succ}\ (\!|y = x + 1|\!)$$

and

$$\models_{tot} (\!|\top|\!)\ \texttt{Succ}\ (\!|y = x + 1|\!)$$

Remark: $\models_{tot} (\!|\Phi|\!)\ P\ (\!|\Psi|\!)$ implies $\models_{par} (\!|\Phi|\!)\ P\ (\!|\Psi|\!)$.

Consider the examples:

Fac2:

```
y = 1;
while (x != 0) {
    y = y * x;
    x = x - 1;
}
```

Sum:

```
z = 0;
while (x > 0) {
    z = z + x;
    x = x - 1;
}
```

The values of y and z are functions of ***the original*** values of x. That value is no longer available as a program variable at the end of the program. We introduce logical variables to handle this situation.

$$\models_{tot} (\!|x = x_0 \wedge x \geq 0|\!) \text{ Fac2 } (\!|y = x_0!|\!)$$

$$\models_{tot} (\!|x = x_0 \wedge x > 0|\!) \text{ sum } \left(\!\left|z = \frac{x_0(x_0 + 1)}{2}\right|\!\right)$$

$$\frac{(\!|\phi|\!)\ C_1\ (\!|\eta|\!) \qquad (\!|\eta|\!)\ C_2\ (\!|\psi|\!)}{(\!|\phi|\!)\ C_1;C_2\ (\!|\psi|\!)}\ \text{Composition}$$

$$\frac{}{(\!|\psi[E/x]|\!)\ x = E\ (\!|\psi|\!)}\ \text{Assignment}$$

$$\frac{(\!|\phi \wedge B|\!)\ C_1\ (\!|\psi|\!) \qquad (\!|\phi \wedge \neg B|\!)\ C_2\ (\!|\psi|\!)}{(\!|\phi|\!)\ \text{if}\ B\ \{C_1\}\ \text{else}\ \{C_2\}\ (\!|\psi|\!)}\ \text{If-statement}$$

$$\frac{(\!|\psi \wedge B|\!)\ C\ (\!|\psi|\!)}{(\!|\psi|\!)\ \text{while}\ B\ \{C\}\ (\!|\psi \wedge \neg B|\!)}\ \text{Partial-while}$$

$$\frac{\vdash \phi' \rightarrow \phi \qquad (\!|\phi|\!)\ C\ (\!|\psi|\!) \qquad \vdash \psi \rightarrow \psi'}{(\!|\phi'|\!)\ C\ (\!|\psi'|\!)}\ \text{Implied}$$

$$\dfrac{\dfrac{(1=1)\; y\; =\; 1\; (y=1)}{(\top)\; y\; =\; 1\; (y=1)}\; i \qquad \dfrac{(y=1 \wedge 0=0)\; z\; =\; 0\; (y=1 \wedge z=0)}{(y=1)\; z\; =\; 0\; (y=1 \wedge z=0)}\; i}{(\top)\; y\; =\; 1;\; z\; =\; 0\; (y=1 \wedge z=0)}\; c$$

$$\cfrac{\cfrac{\left(y \cdot (z+1) = (z+1)!\right) \; \mathtt{z \; = \; z+1} \; \left(y \cdot z = z!\right)}{\left(y = z! \wedge z \neq x\right) \; \mathtt{z \; = \; z+1} \; \left(y \cdot z = z!\right)} \; i \qquad \left(y \cdot z = z!\right) \; \mathtt{y \; = \; y*z} \; \left(y = z!\right)}{\cfrac{\left(y = z! \wedge z \neq x\right) \; \mathtt{z \; = \; z+1; \; y \; = \; y*z} \; \left(y = z!\right)}{\cfrac{\left(y = z!\right) \; \mathtt{while \; (z \; != \; x) \; \{z \; = \; z+1; \; y \; = \; y*z\}} \; \left(y = z! \wedge z = x\right)}{\left(y = 1 \wedge z = 0\right) \; \mathtt{while \; (z \; != \; x) \; \{z \; = \; z+1; \; y \; = \; y*z\}} \; \left(y = x!\right)} \; i} \; w} \; c$$

Using the rule for composition, we get

$$(\top)\ y = 1;\ z = 0;\ \texttt{while}\ (z\ != \ x)\ \{z = z+1;\ y = y*z\}\ (y = x!)$$

The rule for sequential composition suggests a more convenient way of presenting proofs in program logic: ***proof tableaux***. We can think of any program of our core programming language as a sequence.

Corresponding tableau:

$C_1;$
$C_2;$
$\vdots$
C_n

Each of the transitions	
$(\!(\Phi_i)\!) \; C_{i+1} \; (\!(\Phi_{i+1})\!)$	
appeals to one of the proof rules	

$(\!(\Phi_0)\!)$
$C_1;$
  $(\!(\Phi_1)\!)$ justifi cation
$C_2;$
  $(\!(\Phi_2)\!)$ justifi cation
$\vdots$ $(\!(\Phi_{n-1})\!)$ justifi cation
$C_2;$
  $(\!(\Phi_n)\!)$ justifi cation

We show $\vdash_{\mathsf{par}} (\!|y = 5|\!) \; \mathsf{x} \; = \; \mathsf{y} \; + \; 1 \; (\!|x = 6|\!)$:

$$(\!|y = 5|\!)$$
$$(\!|y + 1 = 6|\!) \qquad \text{Implied}$$
$$\mathsf{x} \; = \; \mathsf{y} \; + \; 1$$
$$(\!|x = 6|\!) \qquad \text{Assignment}$$

We prove $\vdash_{\mathsf{par}} (\!|y < 3|\!) \; \mathsf{y} \; = \; \mathsf{y} \; + \; 1 \; (\!|y < 4|\!)$:

$$(\!|y < 3|\!)$$
$$(\!|y + 1 < 4|\!) \qquad \text{Implied}$$
$$\mathsf{y} \; = \; \mathsf{y} \; + \; 1;$$
$$(\!|y < 4|\!) \qquad \text{Assignment}$$

$$(\!|\top|\!)$$
$$(\!|(x + 1 - 1 = 0 \to 1 = x + 1) \land (\neg(x + 1 - 1 = 0) \to x + 1 = x + 1)|\!) \quad \text{Implied}$$

```
a = x + 1;
```
$$(\!|(a - 1 = 0 \to 1 = x + 1) \land (\neg(a - 1 = 0) \to a = x + 1)|\!) \quad \text{Assignment}$$

```
if (a - 1 == 0) {
```
$$\qquad (\!|1 = x + 1|\!) \quad \text{If-Statement}$$

```
    y = 1;
```
$$\qquad (\!|y = x + 1|\!) \quad \text{Assignment}$$

```
} else {
```
$$\qquad (\!|a = x + 1|\!) \quad \text{If-Statement}$$

```
    y = a;
```
$$\qquad (\!|y = x + 1|\!) \quad \text{Assignment}$$

```
}
```
$$\quad (\!|y = x + 1|\!) \quad \text{If-Statement}$$

Definition: An *invariant* of the while-statement `while` B $\{C\}$ having guard B and body C is a formula η such that $\models_{par} (\!| \eta \wedge B |\!) \, C \, (\!| \eta |\!)$; i.e., if η and B are true in a state and C is executed and terminates, then η is again true in the resulting state.

Example:

```
y = 1;
z = 0;
while (z != x) {
    z = z + 1;
    y = y * z;
}
```

iteration	z	y	B
0	0	1	true
1	1	1	true
2	2	2	true
3	3	6	true
4	4	24	true
5	5	120	true
6	6	720	false

Invariant: `y = z!`

Example

$$(\!|\top|\!)$$
$$(\!|1 = 0!|\!)$$ Implied
```
y = 1;
```
$$(\!|y = 0!|\!)$$ Assignment
```
z = 0;
```
$$(\!|y = z!|\!)$$ Assignment
```
while (z != x) {
```
$$(\!|y = z! \wedge z \neq x|\!)$$ Invariant Hyp. $\wedge$ guard
$$(\!|y \cdot (z + 1) = (z + 1)!|\!)$$ Implied
```
    z = z + 1;
```
$$(\!|y \cdot z = z!|\!)$$ Assignment
```
    y = y * z;
```
$$(\!|y = z!|\!)$$ Assignment
```
}
```
$$(\!|y = z! \wedge \neg(z \neq x)|\!)$$ Partial-while
$$(\!|y = x!|\!)$$ Implied

$$\frac{(\!|\eta \wedge B \wedge 0 \leq E = E_0|\!) \, C \, (\!|\eta \wedge 0 \leq E < E_0|\!)}{(\!|\eta \wedge 0 \leq E|\!) \; \texttt{while } B \; \{C\} \; (\!|\eta \wedge \neg B|\!)} \quad \text{Total-while}$$

$$(x \geq 0)$$
$$(1 = 0! \wedge 0 \leq x - 0) \qquad \text{Implied}$$
```
y = 1;
```
$$(y = 0! \wedge 0 \leq x - 0) \qquad \text{Assignment}$$
```
z = 0;
```
$$(y = z! \wedge 0 \leq x - z) \qquad \text{Assignment}$$
```
while (x != z) {
```
$$(y = z! \wedge x \neq z \wedge 0 \leq x - z = E_0) \qquad \text{Invariant Hyp.} \wedge \text{guard}$$
$$(y \cdot (z + 1) = (z + 1)! \wedge 0 \leq x - (z + 1) < E_0) \qquad \text{Implied}$$
```
    z = z + 1;
```
$$(y \cdot z = z! \wedge 0 \leq x - z < E_0) \qquad \text{Assignment}$$
```
    y = y * z;
```
$$(y = z! \wedge 0 \leq x - z < E_0) \qquad \text{Assignment}$$
```
}
```
$$(y = z! \wedge x = z) \qquad \text{Total-while}$$
$$(y = x!) \qquad \text{Implied}$$